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#### A pipelined Reed-Solomon decoder based on a modified step-by-step algorithm

**Key words:** Reed-Solomon codes, Step-by-step algorithm, Ultrawideband, Pipelined structure

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## Motivation

- The most well-known decoding methods for RS codes share three main steps: syndrome calculator (SC), key equation solver (KES), and Chien search and error evaluator (CSEE).
- Massey (1965) presented a totally different decoding algorithm, can directly determine whether the received symbol is erroneous or not and immediately find the corresponding error value without solving the key equations.
- However, Chen *et al.* (2000) and Liu *et al.* (2007) took two different ways to decode when the error number is equal to t or not, which involve a lot of redundant operations.

## Main idea

- A modified step-by-step algorithm is proposed, which merges the two cases and hence reduces the complexity significantly.
- Since the SC block takes a long time, rather than reducing the latency efficiently, the parallel algorithm (Liu *et al.*, 2007) increases the hardware requirements. In this paper, a pipelined architecture is proposed. This architecture does not introduce extra latency while reducing hardware complexity, which increases the hardware utilization efficiency significantly.

#### Method

 A modified step-by-step algorithm is proposed. Algorithm 1 Modified step-by-step algorithm

- 1: Calculate the syndrome values by  $S_i = r(\alpha^i)$  (i = 1, 2, ..., 2t), and simply set  $S_{2t+1} = 0$ .
- 2: Calculate det $(L_k)$  and det $(L_k^{xx})$   $(k = 0, 1, \ldots, t + 1; x = 0, 1, \ldots, k)$ . Then determine the error number v. Thereafter, select and save det $(L_v)$ , det $(L_{v+1})$ , det $(L_{v,0}^{xx})$   $(x = 0, 1, \ldots, t)$ , and det $(L_{v+1,0}^{xx})$   $(x = 0, 1, \ldots, t+1)$ .
- 3: For each symbol  $r_j$   $(j = 0, 1, \ldots, n)$
- 3.1: Calculate  $H_{v,j} = \sum_{x=1}^{t}$  $H_{v+1} = -$

$$H_{v,j} = \sum_{x=1}^{b} \det(L_{v,j}^{xx});$$

$$H_{v+1,j} = \sum_{x=1}^{t+1} \det(L_{v+1,j}^{xx});$$

$$\beta = \det(L_{v})/H_{v,j};$$

$$\det(L_{v,j+1}^{xx}) = \alpha^{(2x-1)} \cdot \det(L_{v,j}^{xx})$$

$$(x = 1, 2, \dots, t);$$

$$\det(L_{v+1,j+1}^{xx}) = \alpha^{(2x-1)} \cdot \det(L_{v+1,j}^{xx})$$

$$(x = 1, 2, \dots, t+1).$$

3.2: Calculate det $(L_{v+1}(\beta, j))$  by det $(L_{v+1}(\beta, j)) = det(L_{v+1}) + \beta \cdot H_{v+1,j}$ . If det $(L_{v+1}(\beta, j)) = 0$ , then  $r'_j = r_j + \beta$ ; otherwise,  $r'_j = r_j$ .

4: Finish

## Method

 2. According to the modified step-by-step algorithm, a pipelined step-by-step decoder is proposed.



Fig. 3 Block diagram of the pipelined step-by-step decoder

# **Major results**

 For RS(23,17) codes, the hardware requirement of the proposed architecture is 42.5%, 24.4%, and 11.3% less than that of ME, pDCME, and the previous SBS architecture, respectively.

| Architecture    | Number<br>of<br>multipliers | Number<br>of<br>adders | Number<br>of<br>multiplexors | Number<br>of<br>registers | Number of<br>constant<br>multipliers | Number<br>of<br>invertors | Total<br>number<br>of XORs |
|-----------------|-----------------------------|------------------------|------------------------------|---------------------------|--------------------------------------|---------------------------|----------------------------|
| ME*             | 25                          | 280                    | 984                          | 2016                      | 13                                   | 1                         | 10090                      |
| pDCME**         | 25                          | 184                    | 488                          | 1408                      | 13                                   | 1                         | 7674                       |
| Previous SBS*** | 45                          | 352                    | 96                           | 400                       | 24                                   | 1                         | 6536                       |
| Proposed        | 38                          | 288                    | 160                          | 424                       | 13                                   | 1                         | 5798                       |

Table 2 Hardware requirement and comparisons of decoders for RS(23,17) codes

\* Lee (2003). \*\* Lee and Lee (2008). \*\*\* Chen *et al.* (2003); Chen and Tasi (2007). ME: modified Euclidean; pDCME: pipelined degree-computationless modified Euclidean; SBS: step-by-step

# **Major results**

 The modified SBS algorithm can always achieve the same performance as the ME and SBS algorithms. That is to say, the modified SBS algorithm can significantly improve the hardware utilization without performance degradation.



Fig. 10 Simulation results for RS codes decoding schemes

#### Conclusions

- A modified step-by-step algorithm and a novel pipelined decoder for RS(23,17) codes are proposed in this paper.
   With modification, the computational complexity is significantly reduced.
- In addition, much less area is needed for this decoder compared with the ME architecture and the pDCME architecture. As a result, the low computational and hardware complexities make the proposed decoder suitable for the UWB system.